Descriptive Set Theory Lecture 3

It's lear from the definitions of as I is that open = Go and closed = For.

Proposition. In a metric your, closed when are GJ. Equily open S. Fo. Proof. let X be a metric space, with metric d. let C be a dosed sub. Let $U_n = \{x \in X : d(x, C) \in \frac{1}{n}\} = \bigcup B(y, \frac{1}{n})$. The cleary, $g \in C$ $f \in \mathcal{A}(u, y)$ $g \in C$ $f \in C$ $f \in \mathcal{A}(u, y)$ $g \in C$ $f \in C$ ffor each up there is xuEC st. d(xu, y) 2 4. Then xy -y here yet been (is dosed.

It's trivial but dosed rubuits of Polich yours are Polish. What other inbests me Polish? E.g. (-a), 1] is Polish being a closed what of IR. What about (0,1]? It turns out that:

There For a Polivil space
$$X_{j}$$
 a subset $Y \le X$ is Polivil if and
only if Y is Gs.
Proof \leftarrow First suppose that $Y = U$ is open. Then the map:
 $G: X \mapsto (X, \overline{d(X, U^{c})})$
is a homeomorphism from U dos a closel subset
of $X \times IR$, dure d is a complete metric for X .
To show the G is continuous, let $X \to X$ all observe
the $G(X_{i}) \to O(k)$. For the continuity of G^{-1} ,
suppose held $G(X_{i}) \to G(k)$, and again observe $X_{i} \to X$.
To show the $G(X) = S(k)$, and again observe $X_{i} \to X$.
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To show the $G(X) = S(k)$ are contracted by $X \to X$.
To show the $G(X) = S(k)$ are completed subset
of $X \times IR^{N}$.

The map X let dx , dy be complete compatible
metrics on X of Y , respectively.
For each u let $U_{i} = he$ union of
all $X \to 0$ with $B \le X$ set.

pe (i) $B \uparrow Y \neq \phi$, (ii) diang (B) < 1, (iii) diandy (BAY) < 1. To show that Y = Mu, fix y EY at show that] Bay satisfying (ii) I (iii). Indeed, since Bdy (y, th) is Y-open, -) an X-open set VEX st. VAY= Bdy (y, th). $ut \quad B := V \cap B_{d_x}(y, \frac{1}{n}).$ To show that Mun = Y, fix x E MUn. Then for each y, Huce is Barx satisfying (i)-(iii). By ii), let yne Bany Since $d(x, y_n) \in dian_x(B_n) \rightarrow O, y_n \xrightarrow{X}$. But by liv), (yn) i, also Candry in dy, so the vo-pleterin of dy gives a limit y EY. Thus, topologically, y x y also in X. But X is Hausdorff here limity are unique, so x=y. Pcop. 2^N is a compact subset of IN^{IN} while IN^{IN} is homeomor-pluic to a Cir subset of 2^{IN} Proof In HWI.

Trees

let A be a nonempty set, which we thick of a set of symbols, an alphabet. A set-merchic tree Ton A is a rubset of A<IN such that T is none-pty is closed down word, i.e. if a word wET then all of initial subwords of w are also int, graph - Theoretic rooted tree from a set - Theoretic one; the converse construction is also easy. For a tree T on A, let [T] := the set of infinite branches through T := Y x & A'N : x | eT}, where if x = (xi) i = 1N, then x | = (X;);<....

Prop. For any tree Ton A the set [T] is closed in A'N, Mare A's given he discrete topology Pool If xn E [T] I xn -> x, then V m E (N, (xn) m E IN) stabilizes al equals x Im, thus x Im ET. Thus, x EST.

Now let YEA'N I define Ty = YwEA''' = JyEY y = w)= = {y| . : y eY al n e IN }. Ty is a tree by definition. In Fact, Ty is a pruned tree, i.e. every wet has a extension, i.e. Jack s.t. wa GT. Also by detimitions, [Ty] = Y.

Prop. A set Y & A'N is loved if al only if Y = [Ty]. In particular, Y -> Ty is a bijection between closed sets in AN al primed trees on A. Proof. <= We already poved but [Ty] is closed. => Suppose Y is closed I let x e [Ty]. Then Vn, ×In ETy. Hence Jy EY s.t. July = ×In. Then Ty y -> x is the ptwise convergence top, so x EY bense Y &, dored. Now let's unlesstand which preved trees yo x correspond to compact sets. König's lenne. Every intruite finity-branching tree I on A has an infimite branch, i.e. [T] = Ø. Proof. Call we Theory if Tw := {v \in I : v > w} is infinite.

Then Q is heavy at since it has only finidely - many Az extensions in T, one of them has to be heavy A lucice using Mt lightmus is timidely-addivine). Keep going ... (technically using the Axion of Dependent Choice). pruned Prop. For a V tree T on A [1] is compact if along if T is finitely - branching, Roof => Suppose Tisu'd findely branching, then FuET st. W hay infinitely many extensions in T. Then [w] := {x = [T] : x > w} is clopen subset w of [T] whit is covered by infinitely may (m) Thus, [u], in it coupard, hence [T] is it congrad. (The sets (ra), we nonempty becase I is proved.) Le let 2 be an open over of IT) and suppose hourds a water diction MA 73 finite subcover. Define an appropriate notion of heaviern for the nodes of T to chow that I x E [T] not covered by U. .. (This is a proof by Jenna Zomback.)